

## Unit 7 Day 3 Section 4.4

- Euler Circuits and Paths

1

### Warm Up ~ Day 3

1. Create a Task Graph based on the Task Table shown.
2. List earliest start times in a table AND mark them on your graph.

Activity	Time	Depends on
A	3	---
B	5	---
C	2	A
D	3	A
E	3	B, D
F	5	C, E
G	1	C
H	2	F, G

Task	EST
A	
B	
C	
D	
E	
F	
G	
H	

3. Minimum Project Time =

4. Critical Path(s) =

5. What is the Latest Start time for D?

6. What is the Latest Start time for C?

2

## Warm Up ~ Day 3

1. Create a Task Graph based on the Task Table shown.
2. List earliest start times in a table AND mark them on your graph.

Activity	Time	Depends on
A	3	----
B	5	----
C	2	A
D	3	A
E	3	B, D
F	5	C, E
G	1	C
H	2	F, G

Task	EST
A	0
B	0
C	3
D	3
E	6
F	6
G	9
H	14

3. Minimum Project Time =

**16 days**

4. Critical Path(s) =

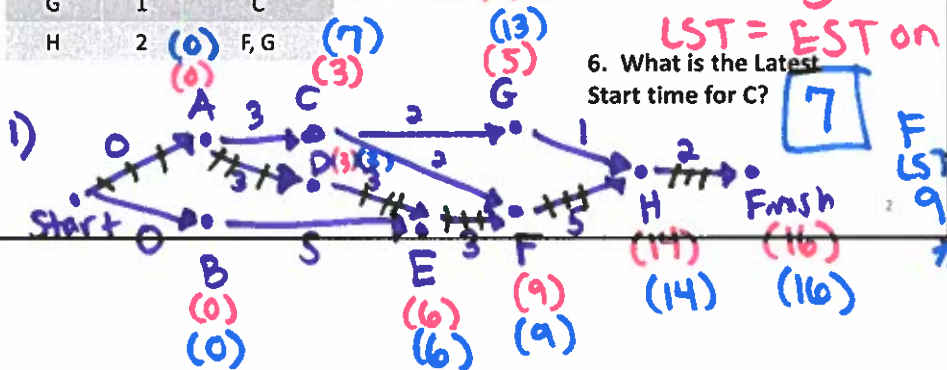
**Start - A - D - E - F - H - Finish**

5. What is the Latest Start time for D?

**3**

6. What is the Latest Start time for C?

**7**



*LST = EST on critical path*

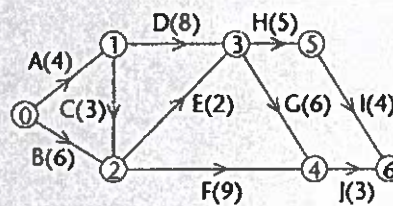
*Time for C = 7 - 2 = 5*

*pick smaller for LST if multiple options*

## Warm Up ~ Day 3 ANSWERS

- List the tasks and earliest start times in a table, as in exercise #1.  
 Determine the minimum project time and all the critical paths.

Determine the critical activities and the length of the critical path for this network.



Activity	Time	Depends on
A	3	----
B	5	----
C	2	A
D	3	A
E	3	B, D
F	5	C, E
G	1	C
H	2	F, G

Task	EST
A	
B	
C	
D	
E	
F	
G	
H	

Minimum Project Time =

Critical Path(s) =

What is the Latest Start time for D?

What is the Latest Start time for C?

*Finish H G C*  
*16 - 2 - 1 = 13*  
*= 11*  
*bigger one does NOT allow time to get other projects on time.*

## Unit 7 Day 3 Section 4.4-4.5

- Euler Circuits and Paths

35

# Euler Circuits and Paths

## Section 4.4



er

**NOT**



ler

36

## Tonight's Homework

In your HW packet...

Packet p. 4 #10

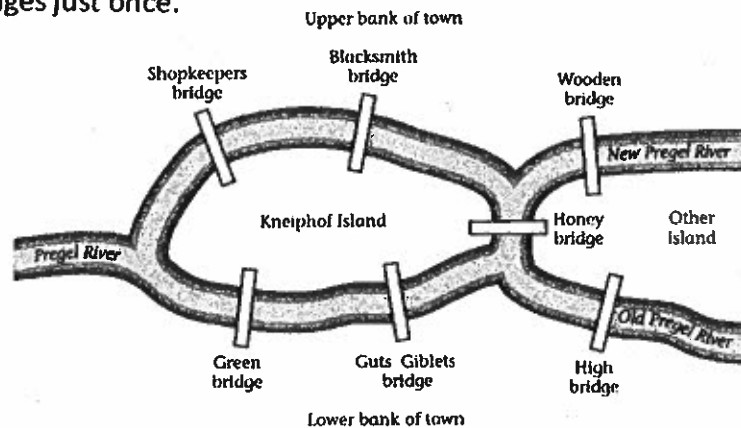
and Packet p. 5

37

37

## The Seven Bridges of Königsberg

The medieval town of Königsberg has a river running through it. There is an island and a fork in the river that together divide the city into four separate land areas. At the time, seven bridges connected the four land areas. The puzzle asked whether it was possible for a stroller to take a walk around the town, crossing each of the seven bridges just once.



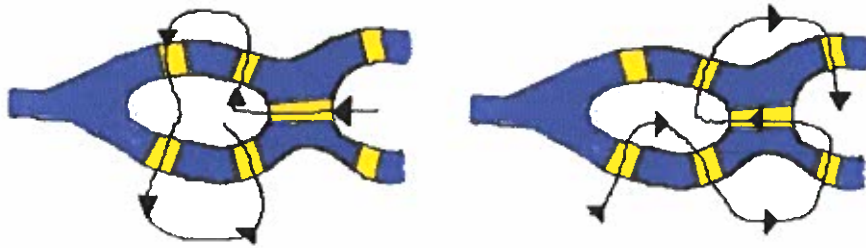
38

## The Seven Bridges of Königsberg

### Solution

- Having trouble? That's okay, so did Euler. It doesn't seem possible to cross every bridge exactly once. In fact it isn't.

### Failed Attempts 😊



39

## Königsberg Problem #2

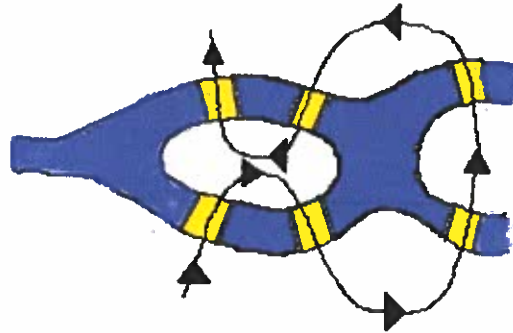
- Suppose they had decided to build one fewer bridge in Königsberg, so that the map looked like this:



40

## Problem #2 Solution

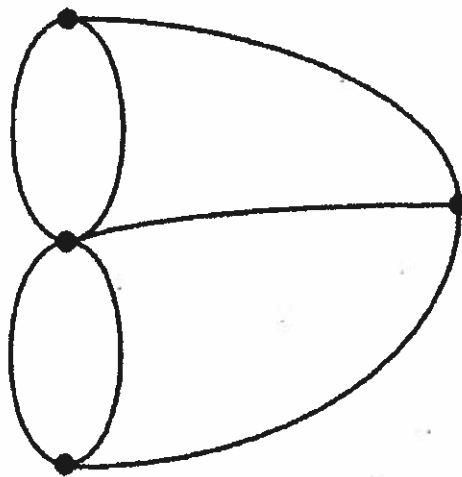
- What makes this one different from the 'real' Konigsberg problem? (Hint: How many bridges lead to each piece of land? Why is having an odd number of bridges leading to a single piece of land problematic?)



41

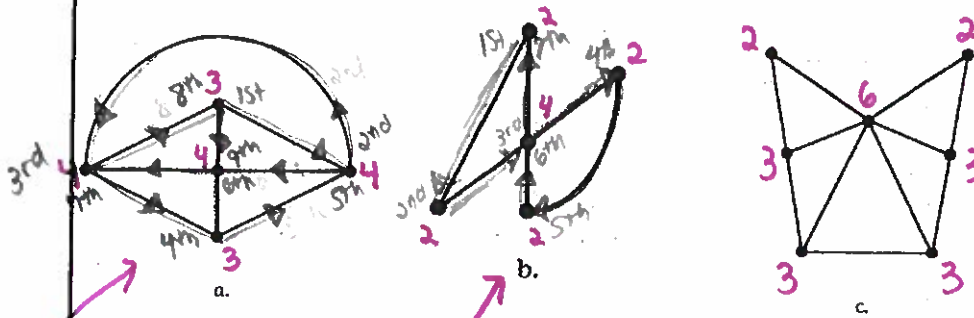
- Need even number of bridge to each piece of land.
- If odd number of bridges to one land area, then example...  
 1<sup>st</sup>) Go to land  
 2<sup>nd</sup>) Go off land  
 3<sup>rd</sup>) Go back on to land  
 AND stuck !!

Attempt to draw this figure without lifting your pencil from the page and without tracing any of the lines more than once.



42

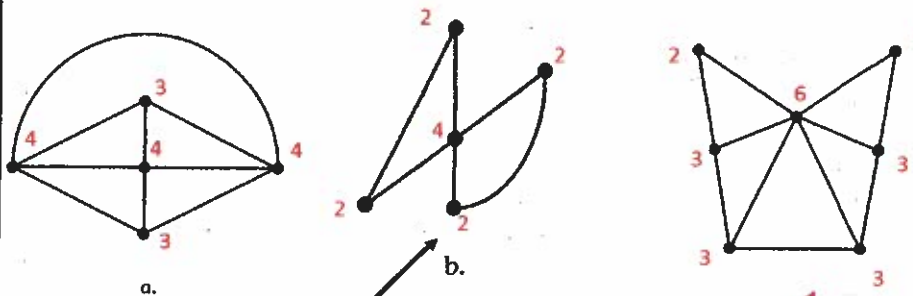
Try to reproduce the following figures without lifting your pencil or tracing the lines more than once. Also, write the degree of each vertex (number of vertices adjacent to it).



1. When can you draw the figures without retracing any edges and still end up at your starting point? *b's figure*
2. When can you draw the figure without retracing and end up at a point other than the one from which you began? *figure a*
3. When can you not draw the figure without retracing? *figure c*

43

**ANSWERS** Try to reproduce the following figures without lifting your pencil or tracing the lines more than once. Also, write the degree of each vertex (number of vertices adjacent to it).



1. When can you draw the figures without retracing any edges and still end up at your starting point?
2. When can you draw the figure without retracing and end up at a point other than the one from which you began?
3. When can you not draw the figure without retracing?

44

Euler: Trace each edge once  
 + don't pick up pencil

**Our Findings (Hopefully)...** \*Write this down!!

You can draw it and you end  
up where you started:

ALL DEGREES ARE EVEN

This is an  
**EULER CIRCUIT**

You can draw it, but you don't  
end where you started:

EXACTLY TWO OF THE  
 VERTICES HAVE AN ODD  
 DEGREE

This is an  
**EULER PATH**

Can't draw it:

NEITHER OF THE  
 PREVIOUS CASES

**Not Possible**

Ex: Figure C  
 from last  
 example

Ex:  
 Figure b  
 from  
 last example

Ex: Figure a  
 from  
 last  
 example

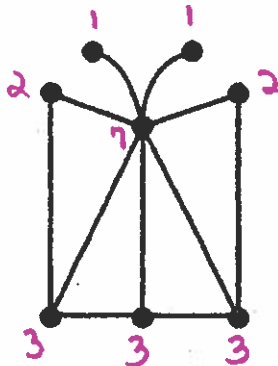
\*Tip:  
 start at odd  
 degree  
 vertex  
 +  
 end at other  
 odd  
 degree  
 vertex

45

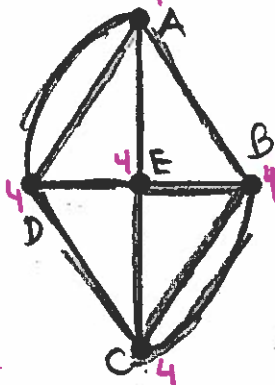
Determine if there is an Euler Circuit, Euler Path, or neither.

If one exists, find a

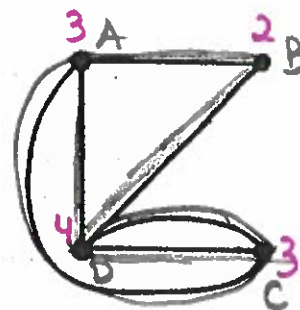
circuit  
 or path.



A Neither  
 (not all even,  
 not 2 odd)



B Euler  
 Circuit  
 (all even degree)



C Euler  
 Path

Euler Circuit: all degrees are even  
 Euler Path: exactly 2 vertices have an odd degree

(exactly 2  
 odd degree)

A B C A D C

Start and end at  
 different ones...  
 at the odd degree  
 vertices

46

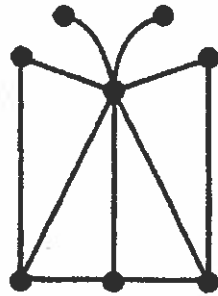
Example: A B C B E D C E A D A  
 start and end are same

\* there are many possible Euler circuits \*

6

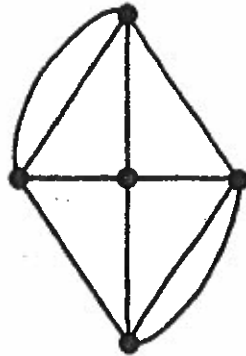


Determine if there is an Euler Circuit, Euler Path, or neither.



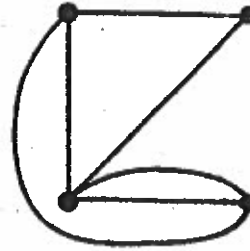
Neither

A



Euler Circuit

B



Euler Path

C

**Euler Circuit:** all degrees are even  
**Euler Path:** exactly 2 vertices have an odd degree

47

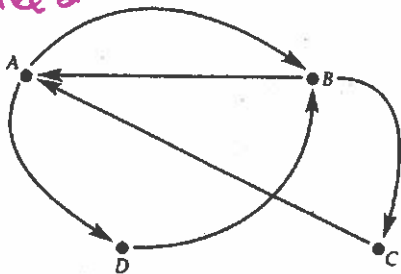
### A "tidbit" on Digraphs

**Directional Graphs** or **DIGRAPHS** are graphs with edges that have direction.

Many applications of graphs require that the edges have direction.

- A city with one-way streets
- A business model with buyers and sellers
- Water flow through a filtration system

A: Indegree 2  
 ↓  
 outdegree 2



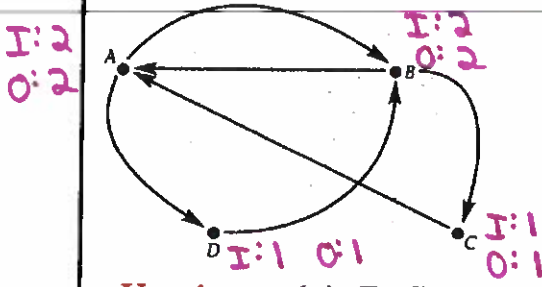
**Indegree** - # of edges coming into a vertex.

**Outdegree** - # of edges going out of a vertex.

D: Indegree 1  
 ↓  
 outdegree 1

48

## A few more "tidbits" on Digraphs



**Indegree** - # of edges coming into a vertex.

**Outdegree** - # of edges going out of a vertex.

**Vertices** - { A, B, C, D }

**Ordered Edges** - { AB, BA, BC, CA, DB, AD }

Notice how the edges are ordered. ... Put starting vertex 1st, then ending one 2nd (one with arrow and)

\* If the **Indegree** and **Outdegree** are equal at every vertex then the Digraph has an Euler Circuit ("Directed" Euler Circuit)

like the diagram above Ex: ABCADBA

If this represented a competition, who would win, A vs C?

A vs D? **A** (arrow starts at A)

**C** (arrow starts at C)

49

you can think of the arrow like a "punch" someone is throwing

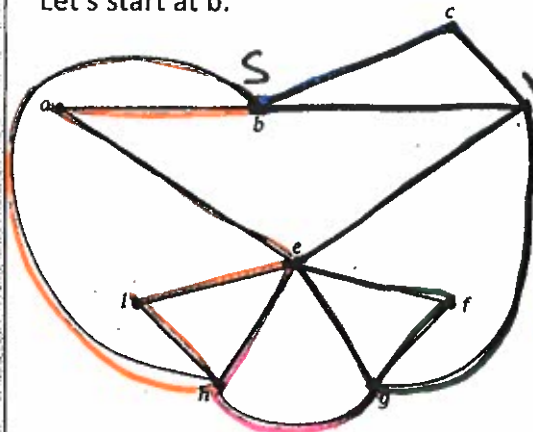
OR C throws punch at A with the arrow pointing at A

When a graph is small, it's easy to find an Euler Circuit by trial & error, but when the graph is bigger you need an algorithm.

### Euler Circuit Algorithm

1. Pick any vertex, and label it S.
2. Construct a circuit, C, that begins and ends at S.
3. If C is a circuit that includes all edges of the graph, go to step 8.
4. Choose a vertex, V, that is in C and has an edge that is not in C.
5. Construct a circuit C' that starts and ends at V using edges not in C.
6. Combine C and C' to form a new circuit. Call this new circuit C.
7. Go to step 3.
8. Stop. C is an Euler circuit for the graph.

Let's start at b.



bcd b  
dgfed  
bcdghfed b  
gheg  
bcd gheg fed b  
hieab h

**Euler Circuit** - cycle with no repetitions of vertices or edges, other than the repetition of the starting and ending vertex

and you must

42

\* TIP! Start with highest degree vertex and/or start with small circuit

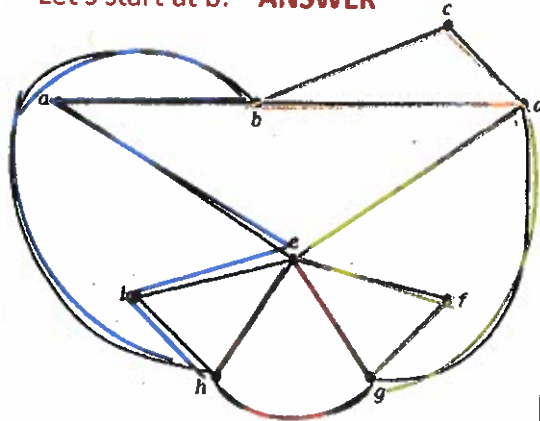
bcdghieabheg fed b

When a graph is small, it's easy to find an Euler Circuit by trial & error, but when the graph is bigger you need an algorithm.

**Euler Circuit Algorithm**

1. Pick any vertex, and label it  $S$ .
2. Construct a circuit,  $C$ , that begins and ends at  $S$ .
3. If  $C$  is a circuit that includes all edges of the graph, go to step 8.
4. Choose a vertex,  $V$ , that is in  $C$  and has an edge that is not in  $C$ .
5. Construct a circuit  $C_1$  that starts and ends at  $V$  using edges not in  $C$ .
6. Combine  $C$  and  $C_1$  to form a new circuit. Call this new circuit  $C$ .
7. Go to step 3.
8. Stop.  $C$  is an Euler circuit for the graph.

Let's start at  $b$ . **ANSWER**



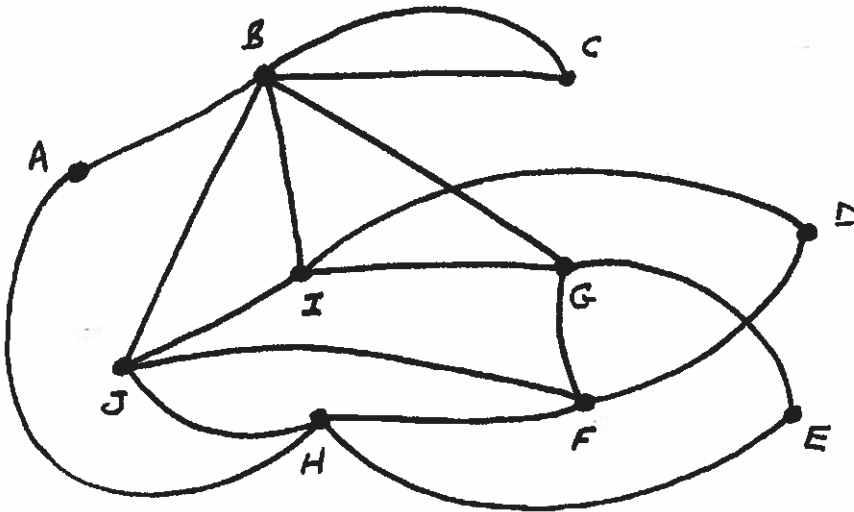
**Circuit** - cycle with no repetitions of vertices or edges, other than the repetition of the starting and ending vertex

51

51

**Euler Circuits Example**

Identify an Euler Circuit by the algorithm.

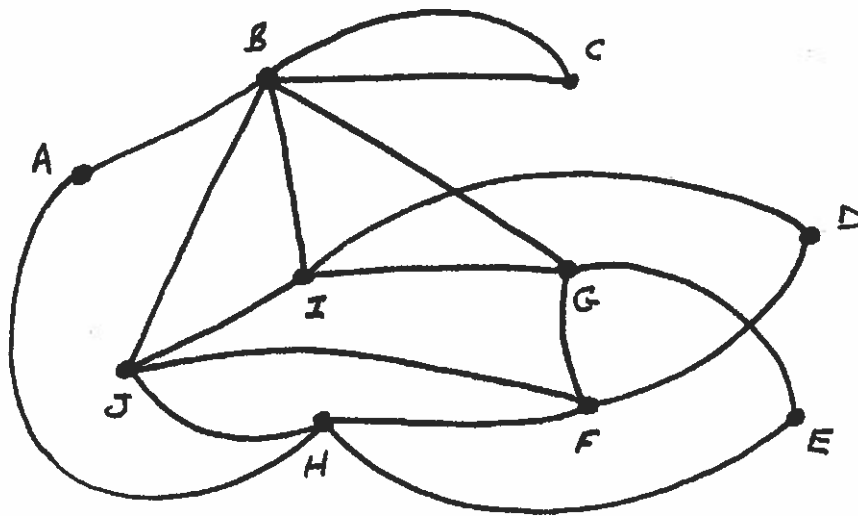


52

52

### Euler Circuits Example ANSWER

Identify an Euler Circuit by the algorithm.



53

### Homework

In your HW packet...

Packet p. 4 #10

and Packet p. 5

54